

4/16/ Disc 2

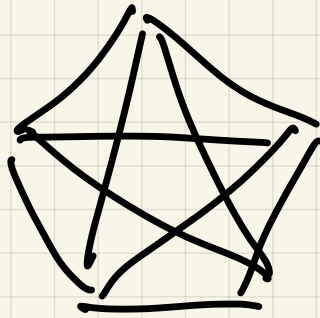
3

Last time  
Simple

Graphs self-  
(no loops  
no // edges)

Special graphs  $K_n, C_n, K_{n,m}$

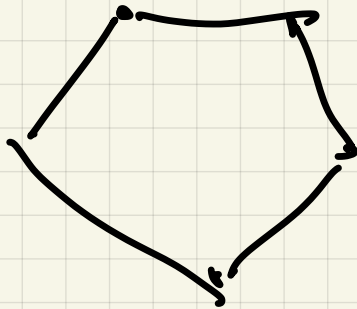
$K_5$



$K_{2,3}$

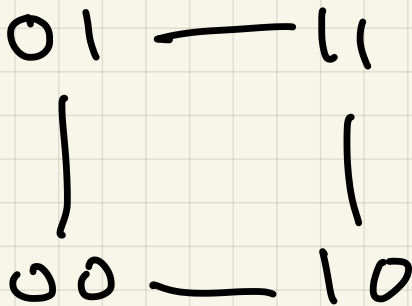


$C_5$



$Q_n =$   
verts n-bit  
strings

$Q_2$



edges if exactly  
one bit  
changes

Adjacency matrix

Adjacency list

Graph isomorphism  $f: G \rightarrow G'$

$f: V_G \rightarrow V_{G'}$  bijection on vertices



s.t.  $(a, b)$  edge in  $G \iff$   
 $(f(a), f(b))$  edge in  $G'$

Notation:  $G \cong G'$

Degree sequence

$G \cong G' \Rightarrow G, G'$  have same  
degree sequence  
same total degree

Ex  $K_3 \cong C_3$



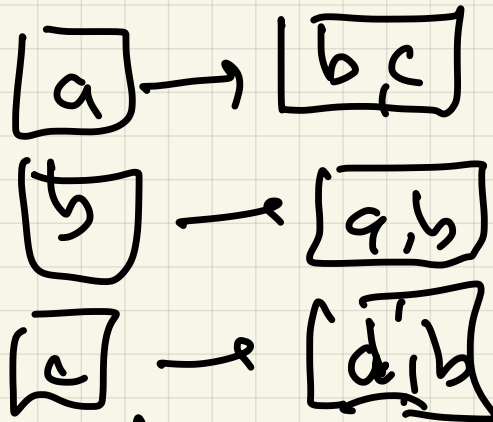
$f(a) = 1$   
 $f(b) = 2$   
 $f(c) = 3$

Degree sequence 2 2 2

Total degree 6

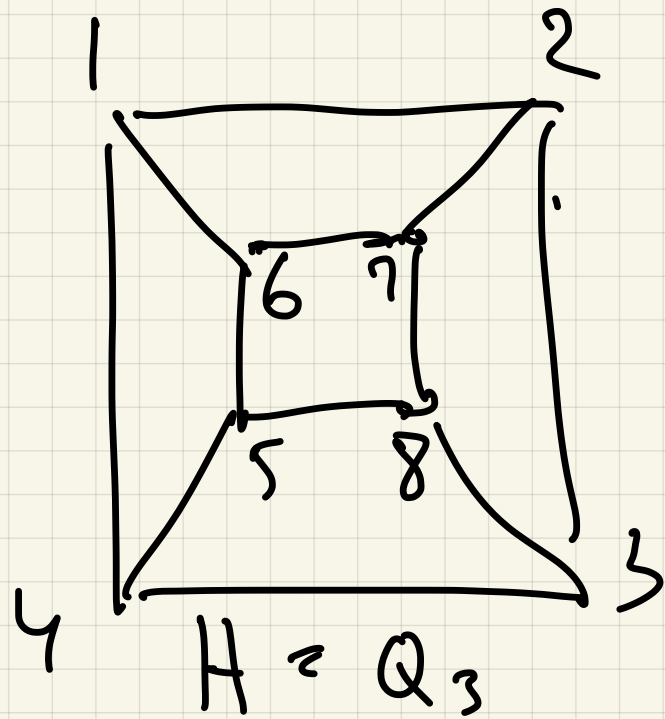
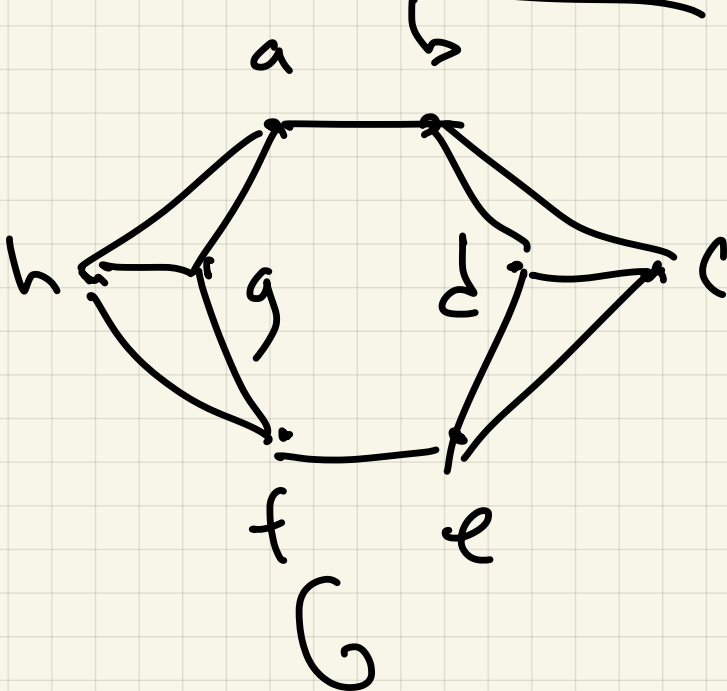
$$\begin{matrix} a \\ b \\ c \end{matrix} \begin{pmatrix} 0 & 1 & 1 \\ 1 & 0 & 1 \\ 1 & 1 & 0 \end{pmatrix}$$

$K_3$  Adj. matrix



Adjacency list

Ex 2  $K_4 \cong G \cong H$



Both 3-regular  
8 vertices

total degree 24

Degree seven 3333333

No:  $G$  has subgraph  $\cong C_3$   
 $H$  does not.

§13.4 Let  $G = (V, E)$

Walk: sequence of vertices

$\langle v_0, v_1, \dots, v_n \rangle$  where

$(v_i, v_{i+1})$  edge

length =  $n$

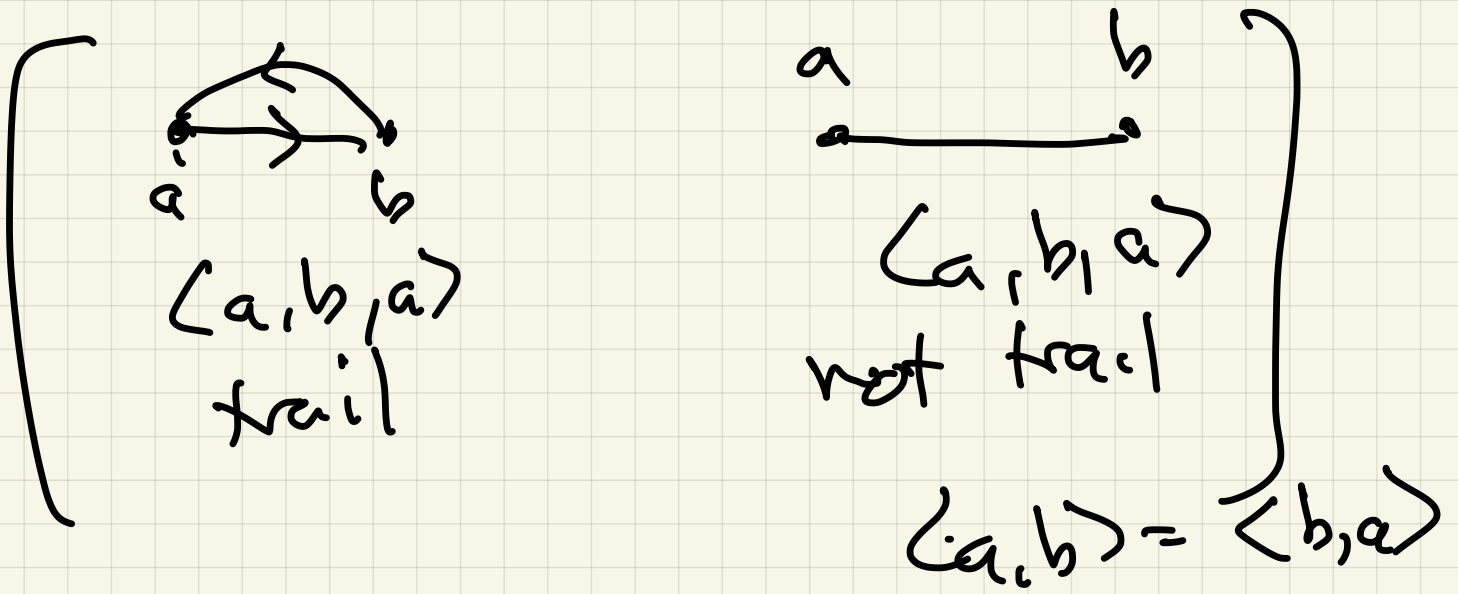
open if  $v_0 \neq v_n$

closed if  $v_0 = v_n$

Types of walks:

Trail

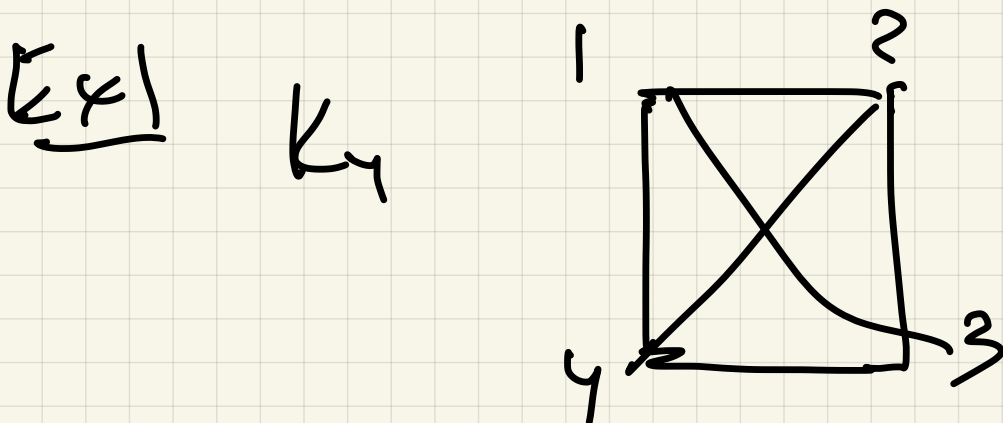
open walk,  
edges different:



Path: Trail with distinct vertices

Circuit: Closed walk with distinct edges

Cycle Circuit with distinct vertices EXCEPT first = last



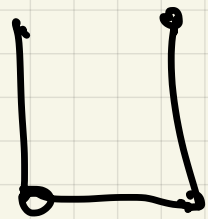
(a)  $\langle 1, 2, 3, 4, 2, 3, 1, 4 \rangle$   
open walk length 7

(b)  $\langle 1, 2, 3, 4, 1, 3 \rangle$   
open walk, trail,  $l=5$ ,  
not path

(c)  $\langle 1, 2, 3, 4, 1 \rangle =$  cycle circuit

longest trail?  $l=5$

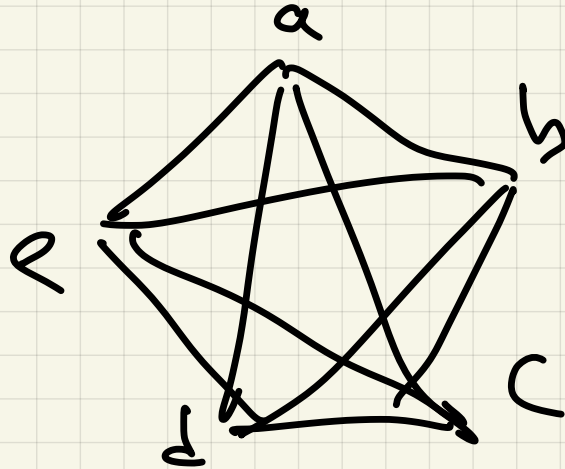
longest path:  $l=3$



longest circuit / cycle  $l=4$

Ex 2

$K_5$



longest cycle  $\ell = 5$

longest path  $\ell = 4$

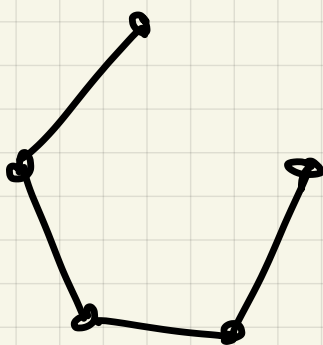
longest circuit?  $\ell = 10$

$\langle a b c d e a d b e c a \rangle$

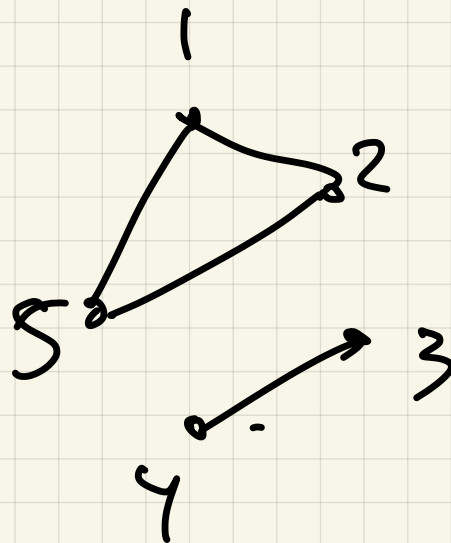
longest trail  $\ell = 9$

$\mathbb{Z} \S 13.5$

G



H



degree sequence is 2 2 2 1 1

but  $G \neq H$ .

G has trail length 4, H doesn't.

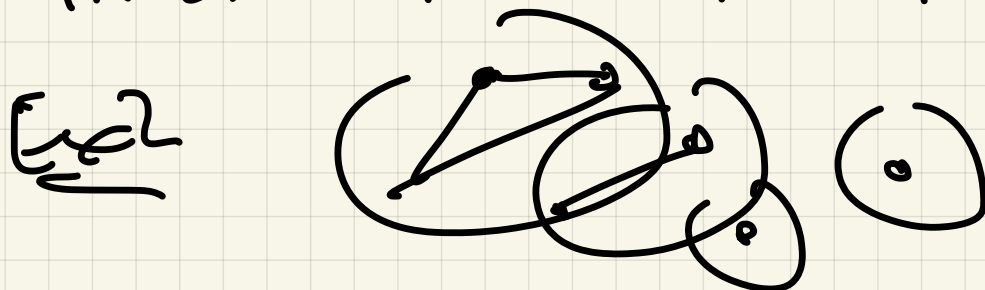
Also,  $G$  connected, it's not.

Definition: Two vertices  $v, w$  in a graph are connected if there's a path from  $v$  to  $w$ .

Graph  $G$  is connected if for each pair  $v, w \in V_G$ ,  $v$  is connected to  $w$ .

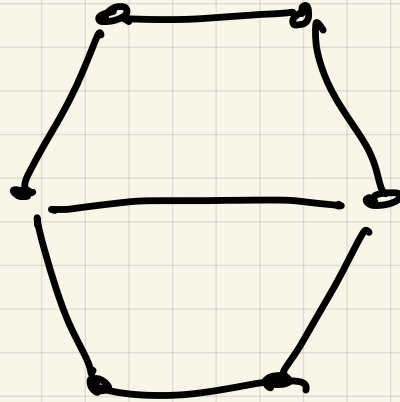
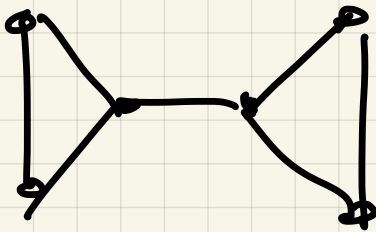
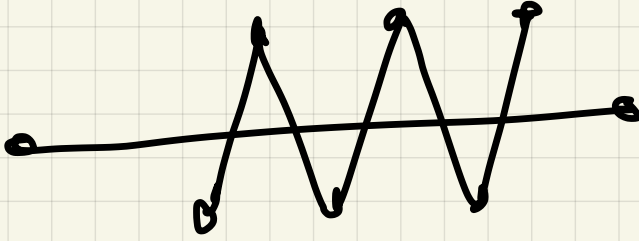
Otherwise  $G$  is disconnected.

A connected component of  $G$  is a maximal set of vertices that are connected.





4 connected components



$G$   
332222

same  
degree sequence

$G \neq H$

removing  
one edge  
disconnects  
 $G$

removing an  
edge  
not disconnects  
 $H$

Connectivity

$G$  is  $k$ -vertex connected if

$G$  has at least  $k+1$  vertices  
and is connected after only  
 $k-1$  vertices are removed

†

Vertex connectivity of  $G$  is

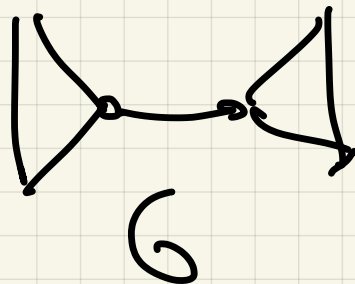
the largest  $k$  with  $G$

$k$ -connected.

Notation

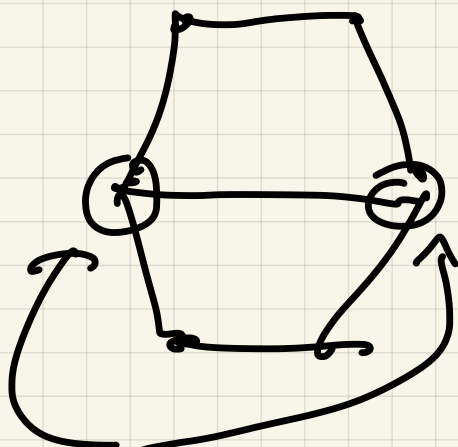
$\kappa(G)$

Ex 1



$$\kappa(G) = 1$$

Ex 2

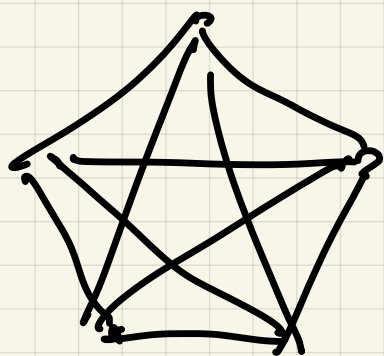


$$\kappa(G) = 2$$

removing these

disconnects the graph

$K_5$



Wuert? by  
definition

$$k(K_n) = n - 1$$

Definition:  $G$  is  $k$ -edge-connected  
if it is connected after  
removing an  $k-1$  edges.

Edge connectivity of  $G$  is

largest such  $k$

||

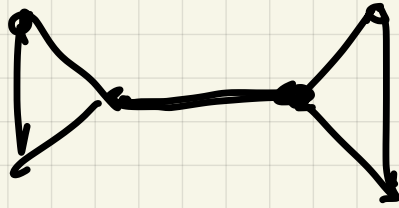
(= min # edges whose removal  
disconnects  $G$ )

Notation  $\lambda(G) =$  edge  
connectivity

Ex

What is  $k(G)$  &  $\lambda(G)$

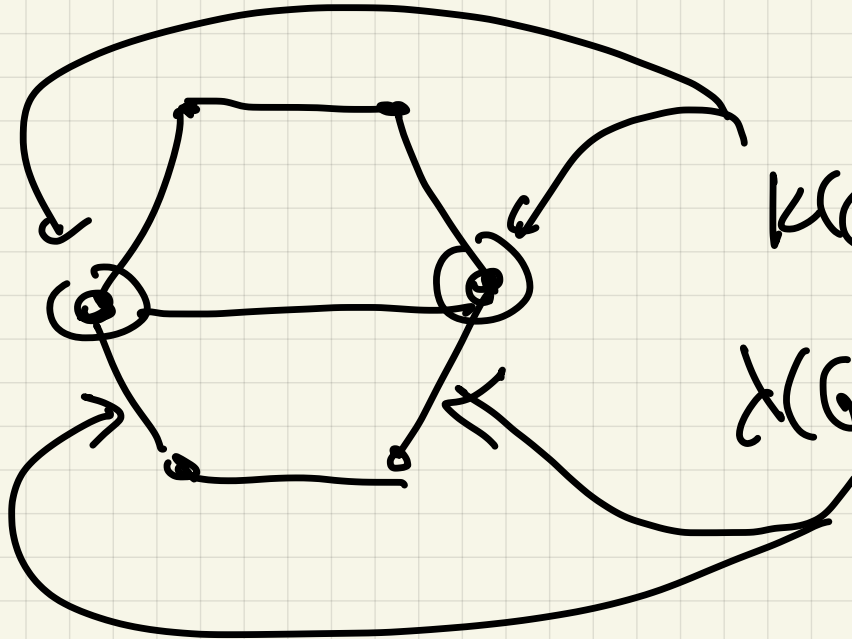
(a)



$$\lambda(G) = 1$$

$$k(G) = 1$$

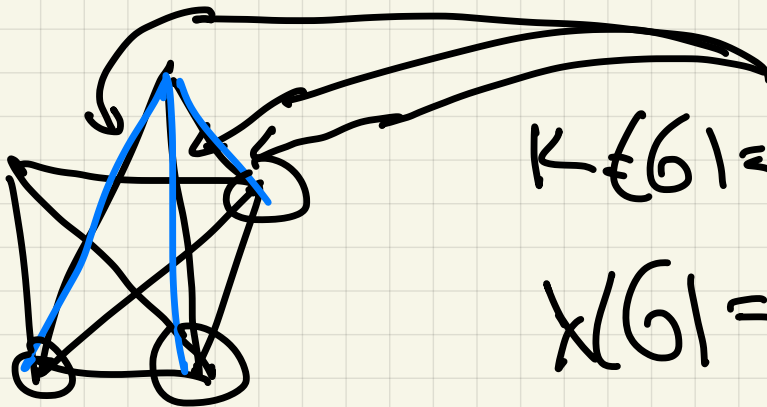
(b)



$$k(G) = 2$$

$$\lambda(G) = 2$$

(c)



$$k(G) = 3$$

$$\lambda(G) = 3$$