

# 9/18/Discrete:

Exam → Sept, 25

Next Wednesday  
Review sheet

Quiz 5

avg 87%

med 90%

$x$	$y$	$x \wedge y$	$y \rightarrow x$	$(y \rightarrow x) \wedge (x \wedge y)$
T	T	T	T	T
T	F	F	T	F
F	T	F	F	F
F	F	F	T	F

$$(x \wedge y) \wedge (y \rightarrow x) \equiv y$$

2.  $y \leftarrow$

Insight: From properties:

De Morgan's  $(x \wedge y) \vee \neg(y \rightarrow x) \stackrel{\text{defn}}{=} (x \wedge y) \vee \neg(\neg y \vee x)$

$$= (x \wedge y) \vee (\neg(\neg y) \wedge \neg x) =$$

$$(x \wedge y) \vee (y \wedge \neg x) =$$

$$(y \wedge x) \vee (y \wedge \neg x) = \text{distrib}$$

$$y \wedge (x \vee \neg x)$$

$$y \wedge T = \textcircled{y}$$

Let's try: Sets

Cardinality

S

|S|

Set builder notation

$$x \in A$$

$$A \subseteq B$$

$$(x \notin A)$$

$$(A \subset B \text{ means})$$

$$A \subseteq B \text{ but}$$

$$A \neq B$$

$$A = B$$

Exo We listed all subsets

of  $A = \{a, b, c, d\}$

$\emptyset = \{\}$ ,  $\{a\}$ ,  $\{b\}$ ,  $\{c\}$ ,  $\{d\}$ ,

$\{a, b\}$ ,  $\{a, c\}$ ,  $\{a, d\}$ ,  $\{b, c\}$ ,  $\{b, d\}$ ,

$\{c, d\}$

$\{a, b, c\}$ ,  $\{a, b, d\}$ ,  $\{a, c, d\}$ ,  $\{b, c, d\}$

$\{a, b, c, d\}$

Defn: If  $A$  is a set,

the power set of  $A$

is the set of all subsets

of  $A$ ,

Notation  $P(A)$  (text =  $2^A$ )

Ex 1  $A = \emptyset$   $P(A) = \{\emptyset\}$

(a)  $|P(A)| = 1$

(b)  $A = \{a\}$ ,  $P(A) = \{\{a\}, \emptyset\}$

(c)  $A = \{a, b\}$

$P(A) = \{\{a, b\}, \{a\}, \{b\}, \emptyset\}$

Pattern  $|P(A)| = 2^{|A|}$

Why?

Here's a way to label the subsets of

$$A = \{a_1, a_2, \dots, a_n\}$$

For  $S \subseteq A$ , I'll label  $S$  by strings of 0s & 1s

$\underline{b_1} \dots \underline{b_n}$

$$b_k = \begin{cases} 0 & a_k \notin S \\ 1 & a_k \in S \end{cases}$$

Ex 1  $A = \{a_1, a_2, a_3\}$

subset	label
$\emptyset$	000
$\{a_1\}$	100
$\{a_2\}$	010
$\{a_3\}$	001
$\{a_1, a_2\}$	110
$\{a_1, a_3\}$	101
$\{a_2, a_3\}$	011
$\{a_1, a_2, a_3\}$	111

There's a 1-1 correspondence

$\left\{ \begin{array}{l} \text{subsets} \\ \text{of } A \end{array} \right\} \longleftrightarrow \left\{ \begin{array}{l} \text{bit strings} \\ \text{of length } n \end{array} \right\}$

↑  
count  
2<sup>n</sup> by  
multiplication  
principle

Ex 2      T/F  
A = {1, 2, 3, 4, 5}

(a)  $2 \in A$       T

(b)  $\{2\} \in A$       F

(c)  $\emptyset \in A$       F

(d)  $\emptyset \subseteq A$       T

(e)  $\emptyset \in P(A)$       T

(f)  $\{2\} \in P(A)$       T

(g)  $\{2\} \in P(A)$       F

(h)  $\{\{2\}\} \in P(A)$       T

$$(i) \{1, \{1, 2\}\} \subseteq A \quad F$$

$$(j) \{1, \{1, 2\}\} \in P(A) \quad F$$

$$(k) \{ \{1\}, \{1, 2\} \} \subseteq P(A) \quad T$$

## § 11 Quantifiers

Defn Existential quantifier

is  $\exists$  it means "there exists"

Ex 1  $a \in \mathbb{Z}$  is odd if

$$(a) \exists c \in \mathbb{Z} \text{ (i) } a = 2c + 1$$

↑  
such that

(b)  $x \in \mathbb{Z}$  is composite if

$$\exists y \in \mathbb{Z} : 1 < y < x \text{ and } \underline{y \mid x}$$

61  $a|b$  if  $\exists c \in \mathbb{Z} : b = ac$

Ex 2  $\exists n \in \mathbb{N} : n^2 = 169$

true :  $n = 13$

$$13^2 = 169 \checkmark$$

Ex 3  $\exists p \in \mathbb{N} : p > 1000$  and  
 $p$  is prime

T/F?

$$p = 1001 = 7 \cdot 11 \cdot 13 \quad \text{not}$$

$$p = 1002$$

$$p = 1003 = 17 \cdot 59$$

$$p = 1004$$

$$1005$$

$p = 1009$  is prime

$\{1, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31\}$   
do not divide 1009



Ex 4  $\exists n \in \mathbb{N}: 2|n$  and  $3|n$   
and  $4|n$  and  $10|n$

proof T: consider  $n = 10!$

smallest such  $n$  is

$$n = 8 \cdot 5 \cdot 9 \cdot 7 = 2620$$

Def The universal quantifier

is  $\forall$  means  
"for every"  
or "for all"

Ex 5  $p \in \mathbb{Z}$  is prime if

$$p > 1 \text{ and } \forall d > 0$$

$$d|p \Rightarrow d=1 \text{ or } d=p$$

Ex 6  $\forall x, y \in \mathbb{Z}$ ,  $x$  odd and  $y$  even  $\Rightarrow$   $x+y$  odd B

Proof: Let  $x$  be odd and  
 $y$  be even.

Then  $\exists c \in \mathbb{Z} : x = 2c + 1$

and  $\exists d \in \mathbb{Z} : y = 2d$

Therefore

$$x + y =$$

$$(2c + 1) + 2d =$$

$\therefore$

$$\frac{2(c + d) + 1}{b/c}$$

$b/c$

$$c + d \in \mathbb{Z} \quad (b/c) \quad c, d \in \mathbb{Z}$$

$\therefore x + y$  is odd.

$\therefore$  therefore

$\therefore$  = because / since

Negating quantifiers

Ex 7 A: There's a man with  
three ears

Symbolically :

A:  $\exists$  man : man has 3 ears

$\neg A$  : There's no man with  
3 ears

also / for each man  $m$ ,  
 $m$  does not have 3 ears

This holds in general:

If  $P(x)$  is a condition on  $x$

[  $P(m)$  :  $m$  has 3 ears ]

$\exists x P(x)$

A

$$\forall x (\neg P(x)) \quad \neg A$$

Punchline:  $\neg(\exists x P(x)) \equiv \forall x (\neg P(x))$

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Similar story with  
universal quantifier

$$\neg(\forall x P(x)) \equiv \exists x (\neg P(x))$$

Ex 8 negate:

$$A: \forall x \in \mathbb{Z} \underbrace{x \text{ is odd}}_{P(x)} \quad F$$

$\neg A$ :

$$\neg(\forall x \in \mathbb{Z} x \text{ odd}) \equiv T$$

$$\exists x \in \mathbb{Z} \neg(x \text{ odd})$$

Ex 9

$\forall x \in \mathbb{Z}$   $x$  is odd or  $x$  is even

$\forall x \in \mathbb{Z}$   $x$  odd  $\vee$   $x$  even (true)

$\neg \forall x \in \mathbb{Z} (x \text{ odd} \vee x \text{ even})$

$\exists x \in \mathbb{Z} \neg (x \text{ odd} \vee x \text{ even})$

$\exists x \in \mathbb{Z} \quad x \text{ not odd} \wedge x \text{ not even}$   
DeMorgan's

$\exists x \in \mathbb{Z} \quad \neg (x \text{ odd}) \wedge \neg (x \text{ even})$